REMARKS

Entry and consideration of this Amendment is respectfully requested.

Respectfully submitted,

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APPENDIX
VERSION WITH MARKINGS TO SHOW CHANGES MADE

IN THE SPECIFICATION:

The specification is changed as follows:

Page 4, first full paragraph:

Each ingress router I_1 of the second domain B uses the SLA information to compute the estimated volume of class-specific traffic between the ingress router I_1 and all egress routers [E_2, E_3]E_1, E_2 in the same domain, to create an $N \times N$ matrix M, where N represents the number of edge routers in the domain. The (i,j)-th element of the traffic matrix for a given class represents the total bandwidth used by that given class from ingress router i to egress router j. For example, as illustrated in Figure 1, for the second domain B, element (1,1) of the matrix M equals λ_1, and element (1,2) of the matrix equals λ_2. Once constructed, the traffic matrices are used to compute the provisioning routes (e.g., paths), for each non-zero element of those matrices, and the computed paths are pinned down using multi-protocol label switching (MPLS) for Diffserv networks or multi-protocol lambda switching (MPλS) for optical networks.

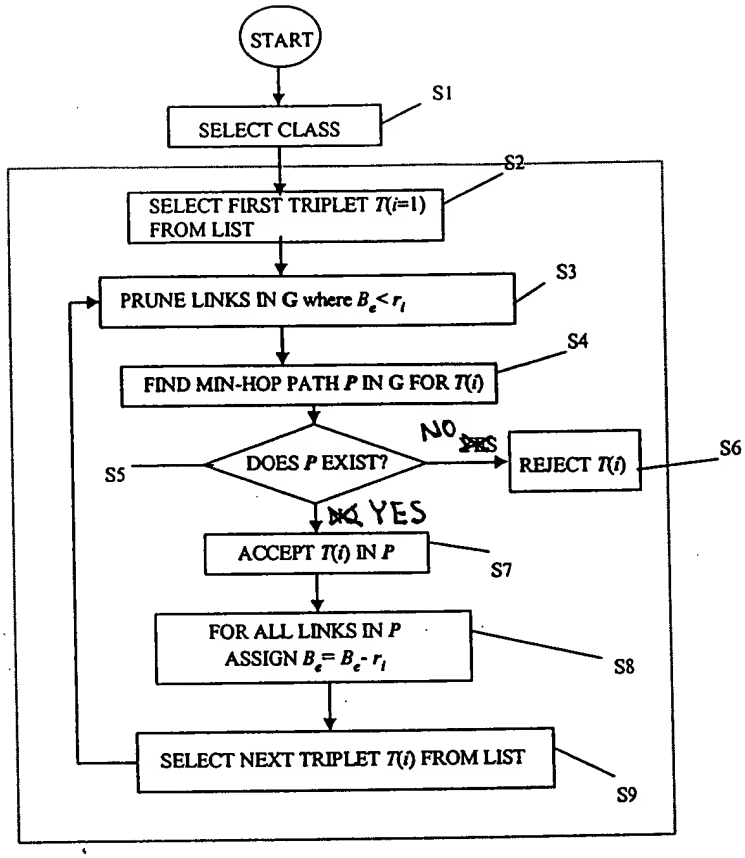
Page 6, after equation (4), please insert the following paragraph:

--V represents the total amount of bandwidth of accepted flows, and W represents the total amount of bandwidth of all flows.--

Page 16, second full paragraph:

Accordingly, in the next step S21, M is defined as the subset of those already accepted (i.e., during the previous i–1 steps) quadruplets $T(1),\ldots, T(i–1)$ for which the following two

conditions hold true. First, the bit $[r_j]\underline{b_j}$ of quadruplet is TRUE and the path SPA(j) thus can be altered. Second, all links e in Q belong to the path SPA(j): $Q \subset \mathrm{SPA}(j)$. Therefore, if the bandwidth reservation for r_j of the quadruplet T(j) for its path SPA(j) is removed, the available bandwidth at each link e in Q increases by r_j. Since the i^{th} flow requires bandwidth reservation of $r_i \leq r_j$, this increase is sufficient for accommodating the i^{th} flow using its path SPI(i).



PRIOR ART

Figure 2

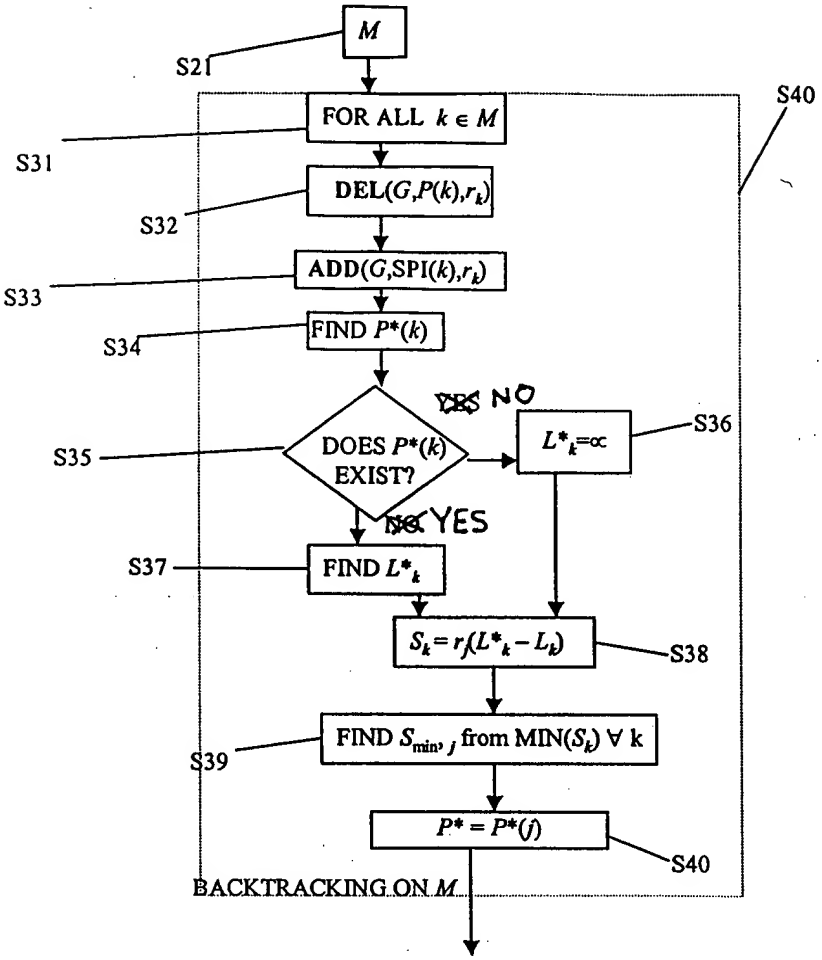

M
S21
FOR ALL $k \in M$
S31
S40
DEL(G,$P(k)$,r_k)
S32
ADD(G,SPI(k),r_k)
S33
FIND $P^*(k)$
S34
NO
DOES $P^*(k)$ EXIST?
S35
$L^*_k=\infty$
S36
YES
FIND L^*_k
S37
$S_k = r_j(L^*_k - L_k)$
S38
FIND $S_{min,j}$ from MIN(S_k) $\forall$ k
S39
$P^* = P^*(j)$
S40
BACKTRACKING ON M

Figure 6